

# Network models

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НАЦИОНАЛЬНЫЙ ИССЛЕДОВАТЕЛЬСКИЙ  
УНИВЕРСИТЕТ

- Power law degree distribution
- Small average distance (diameter)
- Large clustering coefficient
- Hierarchical structure (often)

## Models:

- Random graphs (Erdos & Renyi, 1959)
- Small world networks (Watts & Strogatz, 1998)
- Preferential attachment (Barabasi & Albert, 1999)
- .....
- Strategic network formation (Jackson & Wolinsky, 1996)

Erdos and Renyi, 1959.

Graph  $G\{E, V\}$ , nodes  $n = |V|$ , edges  $m = |E|$

Random graph models

- $G_{n,m}$ , a randomly selected graph from the set of  $C_{n(n-1)/2}^m$  graphs with  $n$  nodes and  $m$  edges
- $G_{n,p}$ , each pair from  $n$  nodes is connected with probability  $p$ ,  
 $m$  - random number
- A graph with  $m$  edges and  $n$  nodes

$$\langle m \rangle = p \frac{n(n-1)}{2}$$

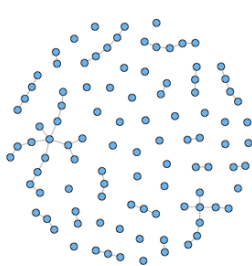
$$\langle k \rangle = \frac{1}{n} \sum_i k_i = \frac{2m}{n} = p(n-1) \approx pn$$

- Graph  $G_{n,p}$  -function of  $p$ :  $p = 0$  - empty graph,  $p = 1$  - full graph
- There are exist critical  $p_c$ , structural changes from  $p < p_c$  to  $p > p_c$
- Critical value:

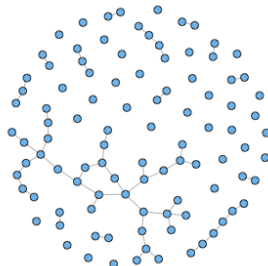
$$p_c = \frac{1}{n}, \quad \langle k \rangle = p_c n = 1$$

- when  $p < p_c$ , ( $\langle k \rangle = pn < 1$ ) there is no components with more than  $O(\ln n)$  nodes, all components are either trees or have maximum one cycle, largest component is a tree
  - when  $p = p_c$ , ( $\langle k \rangle = pn = 1$ ) the largest component has  $O(n^{2/3})$  nodes
  - when  $p > p_c$ , ( $\langle k \rangle = pn > 1$ ) gigantic component has all  $O(n)$  nodes
- Phase transition:  
separate connected islands  $\rightarrow$  gigantic connected component

# Random graph model

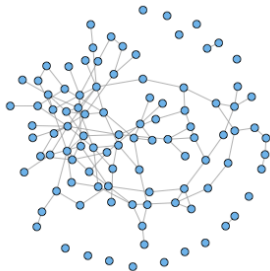


$$p < p_c$$

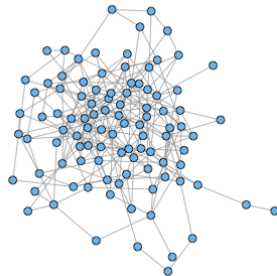


$$p = p_c$$

# Random graph model



$$p > p_c$$



$$p \gg p_c$$

R: `erdos.renyi.game` `{igraph}`

- Node degree distribution - Poisson ( $n \rightarrow \infty$  at fixed  $\langle k \rangle = pn$ ):

$$P(k) = \frac{\langle k \rangle^k e^{-\langle k \rangle}}{k!}$$

- Graph diameter when  $p > p_c$ :

$$d = \frac{\ln n}{\ln \langle k \rangle}$$

- Clustering coefficient:

$$C = p = \frac{\langle k \rangle}{n}$$

$$C \rightarrow 0 \text{ when } n \rightarrow \infty$$

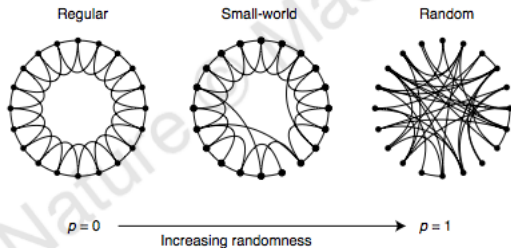


Watts and Strogatz, 1998

One parameter model, interpolation between regular lattice and random graph

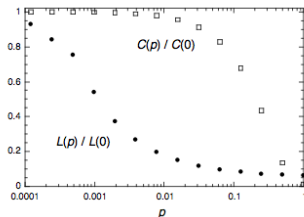
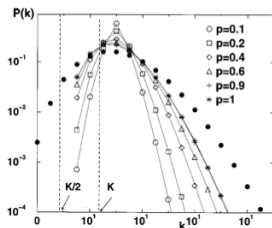
- start with regular lattice with  $n$  - nodes,  $k$  - nearest neighbors (node degree)
- randomly connect with other nodes with probability  $p$ , forms  $pnk/2$  "long distance" connections
- $p = 0$  - regular lattice,  $p = 1$  - random graph

# Small world model

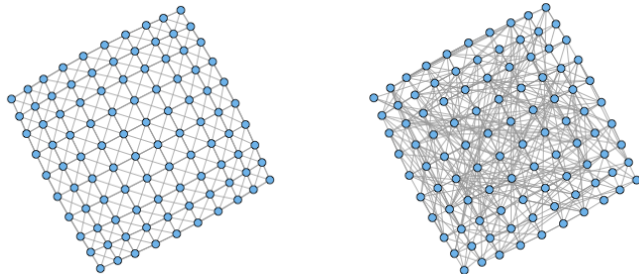


# Small world model

- Node degree distribution:  
Poisson like
- Ave. path length  $\langle L(p) \rangle$  :  
 $p \rightarrow 0$ , ring lattice,  $\langle L(0) \rangle = 2n/k$   
 $p \rightarrow 1$ , random graph,  $\langle L(1) \rangle = \log(n)/\log(k)$
- Clustering coefficient  $C(p)$  :  
 $p \rightarrow 0$ , ring lattice,  $C(0) = 3/4 = \text{const}$   
 $p \rightarrow 1$ , random graph,  $C(1) = k/n$



# Small world model



20% rewiring:

ave. path length = 3.58  $\rightarrow$  ave. path length = 2.32

clust. coeff = 0.49  $\rightarrow$  clust. coeff = 0.19

R: `watts.strogatz.game`, `graph.lattice,rewire.edges` {igraph}

# Preferential attachment model

Barabasi and Albert, 1999

Dynamical (growth) model

- $t = 0$ ,  $n_0$  unconnected nodes
- growth: on every step add a node with  $m$  edges ( $m \leq n_0$ )
- Preferential attachment: probability of linking to existing node is proportional to the node degree  $k_i$

$$k_i(t + \delta t) = k_i(t) + \frac{k_i(t)}{2t} \delta t$$

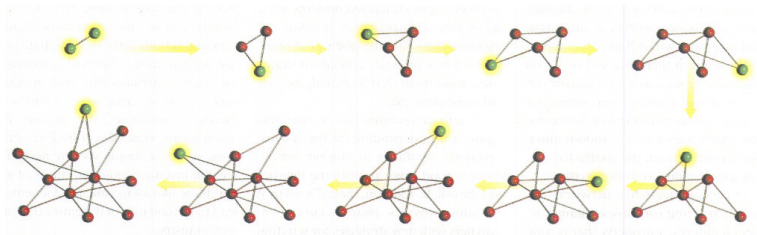
$$\frac{dk_i(t)}{dt} = \frac{k_i(t)}{2t}$$

initial conditions:  $k_i(t_i) = m$

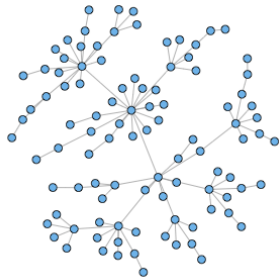
- Time evolution of a node degree

$$k_i(t) = m \left( \frac{t}{t_i} \right)^{1/2}$$

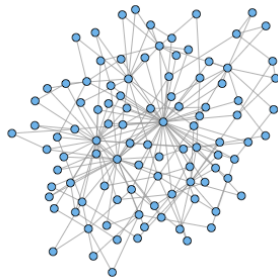
# Preferential attachment model



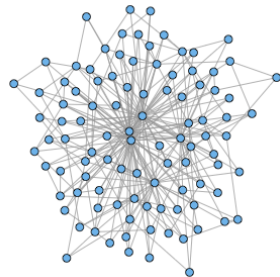
# Preferential attachment model



$m = 1$



$m = 2$



$m = 3$

R: `barabasi.game` {igraph}

# Preferential attachment model

- Node distribution function - "power law":

$$P(k) = \frac{2m^2}{k^3}$$

- Average path length:

$$\langle L \rangle = \log(n) / \log(\log(n))$$

- Clustering coefficient:

$$C \sim \frac{1}{n^{0.75}}$$



# Model comparison

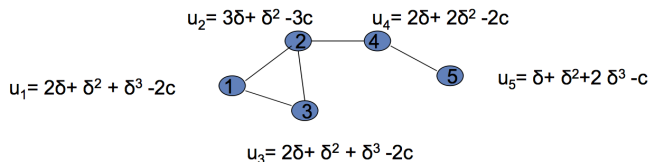
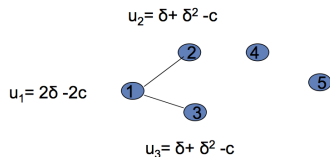
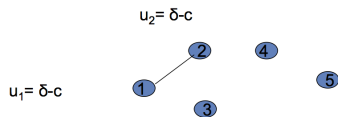
	Random	BA model	WS model	Real World
$P(k)$	$\frac{z^k e^{-z}}{k!}$	$k^{-3}$	poisson like	power law
$C$	$p = \frac{\langle k \rangle}{n}$	$n^{-0.75}$	const	large
$\langle L \rangle$	$\frac{\log(n)}{\log(\langle k \rangle)}$	$\frac{\log(n)}{\log \log(n)}$	$\log(n)$	small world

Jackson and Wolinsky, 1996

"A Strategic Model of Social and Economic Networks"

- why networks become the way they are
- people (agents) making rational choices establishing connections
- maximizing individual utility (incentives)
- connections bring costs and benefits
- stability of the network
- social efficiency (best for the society)
- friendship, professional, political, trade networks

# Distance-based utility function



- Distance-based utility

$$u_i(G) = \sum_{j \in N_i^{n-1}} \delta^{l_{ij}} - d_i \cdot c$$

$l_{ij}$  - shortest path,  $d_i$  - node degree,  $\delta$ ,  $C$  - parameters,  $\delta < 1$

- Co-author model

$$u_i(G) = \sum_{j \in N_i} \left( \frac{1}{d_i} + \frac{1}{d_j} + \frac{1}{d_i d_j} \right), \quad d_i \neq 0$$

$u_i(G) = 1$  if  $d_i = 0$ ,  $d_i$  - node degree

# Network stability and efficiency

- Evolution: forming a link - mutual consent, removing a link - one person decision
- Network is pairwise stable if no player wants to remove a link and no two players want to add a link:

$$\forall i \ u_i(G) \geq u_i(G - e_{ij})$$

$$\forall i, j \text{ if } u_i(G + e_{ij}) > u_i(G), \text{ then } u_j(G + e_{ij}) < u_j(G)$$

- Strong efficiency ("best network", maximizing total utility for the society):

$$G^* = \max_G \sum_i u_i(G)$$

- Pareto efficiency (impossible to increase anyone's utility without decreasing someone's else):

$$\nexists G' : \begin{aligned} u_i(G') &\geq u_i(G) \text{ for all } i \\ \text{and } u_i(G') &> u_i(G) \text{ for one } i \end{aligned}$$

- Complex networks (degree distr, diameter, clustering coeff)
- Node and link analysis (centrality, pagerank)
- Network communities (modularity, detection methods)
- Node similarity (structural equiv, distance, similarity matrix)
- Network structure (cliques, cores, motifs, dyads, triads)
- Network visualization
- Network models (random, small world, pref. attachment)